

# CHAPTER 8: EQUILIBRIUM WITH COMPLETE MARKETS

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Equilibrium in pure exchange, infinite horizon economies, with agents facing stochastic endowments.

Two market structures:

- 1 Arrow-Debreu structure—complete markets in dated contingent claims (time  $\tau$ , state vector  $s^\tau$ ), all traded at time 0.
- 2 Sequential, complete, market structure: Arrow securities traded one period in advance; trades occur at each  $t \geq 0$ .

These two market structures result in identical consumption allocations.

# MARKET SETTING

At  $t \geq 0$ , a stochastic event  $s_t$  is drawn from the state space  $S$ .

Denote the history of events up to and inclusive the one at  $t$  as  $s^t = (s_0, s_1, \dots, s_t)$ , with the *unconditional* probability of observing  $s^t$ ,  $\pi_t(s^t)$ .

The *conditional* probability of observing  $s^t$ , given  $s^\tau$  was realized, is  $\pi_t(s^t | s^\tau)$ ,  $t \geq \tau$ .

$s^t$  is publicly observable.

Assume  $\pi_0(s_0) = 1$ ,  $s_0$  given.

## MARKET SETTING—CONTD.

- Agents:  $i = 1, 2, \dots, I$ .
- (Stochastic) Endowments of agent  $i$  at  $t$ :  $y_t^i(s^t)$ .
- At time 0, household  $i$  purchases a history dependent consumption plan,  $c^i = \{c_t^i(s^t)\}_{t=0}^{\infty}$ . Ranks consumption streams with:

$$U^i = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t u [c_t^i(s^t)] \pi_t(s^t).$$

- $u—C^2$ , strictly concave,  $c \geq 0$ .
- A feasible allocation must satisfy:

$$\sum_i c_t^i(s^t) \leq \sum_i y_t^i(s^t), \text{ all } t, s^t.$$

# EQUILIBRIUM OUTCOMES AND HISTORY DEPENDENCE

Individual  $i$ 's endowment history up to time  $t$  can be described as:  $y_0^i(s_0), y_1^i(s^1), y_2^i(s^2), \dots, y_t^i(s^t)$ . (Note the general dependence of the time  $t$  endowment realization on the entire history of stochastic process, preceding time  $t$ .)

- In the complete markets models, consumption allocation *does not* depend on individual history, and is *only* a function of the aggregate endowment.

# PARETO PROBLEM

Solve the planner's problem:

$$W = \max \sum_{i=1}^l \lambda_i U(c^i) \quad (8.4.1)$$

$$\text{s.t. } \sum_i c_t^i(s^t) \leq \sum_i y_t^i(s^t), \quad (8.2.2)$$

where the Pareto weights  $\lambda_i \geq 0$ ,  $i = 1, \dots, l$ .

## PARETO PROBLEM—CONTD.

Assign non-negative Lagrange multipliers  $\theta_t(s^t)$  for each  $t$  and  $s^t$  constraint (8.2.2), and form the planner's Lagrangian:

$$L = \sum_{t=0} \sum_{s^t} \left[ \sum_i \lambda_i \beta^t \pi_t(s^t) u [c_t^i(s^t)] + \theta_t(s^t) \sum_i [y_t^i(s^t) - c_t^i(s^t)] \right]$$

For any  $i$ , the FOC with respect to  $c_t^i(s^t)$  is:

$$\beta^t u' [c_t^i(s^t)] \pi_t(s^t) = \lambda_i^{-1} \theta_t(s^t)$$

## EFFICIENT ALLOCATION

Thus, the ratio of the FOCs for  $i$  and 1 is:

$$\frac{u'(c_t^i(s^t))}{u'(c_t^1(s^t))} = \frac{\lambda_1}{\lambda_i},$$

$$u'(c_t^i(s^t)) = \lambda_i^{-1} \lambda_1 u'(c_t^1(s^t))$$

$$c_t^i(s^t) = (u')^{-1} [\lambda_i^{-1} \lambda_1 u'(c_t^1(s^t))] \quad (8.4.3)$$

$$\sum_i c_t^i(s^t) = \sum_i (u')^{-1} [\lambda_i^{-1} \lambda_1 u'(c_t^1(s^t))] = \sum_i y_t^i(s^t).$$

Thus,  $c_t^1(s^t)$  is a function of the aggregate endowment *only*, and *not* dependent on the specific history  $s^t$  leading to that outcome, *nor* on the realization of individual endowments.

That is,

$$c_t^i(s^t) = c_\tau^i(s^\tau) \quad \text{when} \quad \sum_i y_t^i(s^t) = \sum_i y_\tau^i(s^\tau).$$

# TIME 0 TRADING: ARROW-DEBREU SECURITIES

- Arrow-Debreu securities: claims to consumption for each possible date  $t$ , history  $s^t$ .
- What is the optimal allocation in a competitive economy with Arrow-Debreu securities?
- Assume: a complete set of state-contingent securities traded at time 0.
- The price of a security at time 0 (**now**) for 1 unit of consumption **to be delivered at  $t$  if  $s^t$  happens** is  $q_t^0(s^t)$ .

# HOUSEHOLD PROBLEM

Household chooses  $c^i(s^t)$  to

$$\max U(c^i) = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t u [c_t^i(s^t)] \pi_t(s^t) \quad (8.2.1)$$

$$\text{s.t.} \quad \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \leq \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t). \quad (8.5.1)$$

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \leq \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t).$$

- No expectations: the household is making all of its lifetime trades at time 0 at a well-defined price vector for all Arrow-Debreu commodities.
- Household buys claims to different “contingent” consumption bundles: delivered if  $s^t$  is realized.

Lagrangian for household  $i$  is:

$$L = \sum_{t=0}^{\infty} \sum_{s^t} \beta^t u [c_t^i(s^t)] \pi_t(s^t) + \mu_i \left[ \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) y_t^i(s^t) - \sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) c_t^i(s^t) \right].$$

F.O.C.:

$$\frac{\partial U(c^i)}{\partial c_t^i(s^t)} = \mu_i q_t^0(s^t),$$

or  $\beta^t u' [c_t^i(s^t)] \pi_t(s^t) = \mu_i q_t^0(s^t)$  (8.5.4)

# COMPETITIVE EQUILIBRIUM

DEFINITIONS. A **price system** is a sequence of functions  $\{q_t^0(s^t)\}_{t=0}^{\infty}$ . An **allocation** is a list of sequences of functions  $c^i = \{c_t^i(s^t)\}_{t=0}^{\infty}$ , one sequence for each  $i$ .

DEFINITION. A **competitive equilibrium** is a *feasible* allocation such that, given the price system, the allocation solves each household's problem.

## EQUILIBRIUM ALLOCATIONS

Equation (8.5.4) implies for any  $i$  and  $j$ :  $\frac{u' [c_t^i(s^t)]}{u' [c_t^j(s^t)]} = \frac{\mu_i}{\mu_j}$  (8.5.5).

That is, the ratios of marginal utilities for any consumers  $i$  and  $j$  are constant across all dates and state histories.

An equilibrium allocation solves (8.2.2), (8.5.1), and (8.5.5).

(8.5.5) implies for any two consumers  $i$  and 1:

$$c_t^i(s^t) = (u')^{-1} \left[ u' (c_t^1(s^t)) \frac{\mu_i}{\mu_1} \right].$$

Summing this expression over  $i$  and utilizing (8.2.2) at equality, we obtain:

$$\sum_i (u')^{-1} \left[ u' [c_t^1(s^t)] \frac{\mu_i}{\mu_1} \right] = \sum_i c_t^i(s^t) = \sum_i y_t^i(s^t).$$

# COMPETITIVE EQUILIBRIUM

The competitive equilibrium allocation is a function of the realized aggregate endowment and depends neither on the specific history leading to that outcome nor on the realizations of individual endowments:  $c_t^i(s^t) = c_\tau^i(s^\tau)$  for all  $s^t$  and  $s^\tau$  for which  $\sum_i y_t^i(s^t) = \sum_i y_\tau^i(s^\tau)$ .

Units of the price system are arbitrary. Let  $q_0^0(s^0) = 1$ ; the rest of the prices will be quoted in the units of time 0 goods.

A competitive equilibrium allocation is a *particular* Pareto optimal allocation, the one for which:

- $\lambda_i = \mu_i^{-1}$ .
- The shadow prices  $\theta_t(s^t) = q_t^0(s^t)$ .

# SEQUENTIAL TRADING OF ARROW SECURITIES

It can be shown that one-period Arrow securities are enough to make the markets complete if the markets are reopened each day, and the trades on *one-period ahead state-contingent* consumption claims occur. This sequential trading attains the same allocation as the competitive equilibrium in the economy with time zero trading of state-contingent claims.

We can price different (redundant) assets using the prices for contingent claims from our competitive complete markets equilibrium.

## EXAMPLE 8.6.1. RISK SHARING WITH CRRA PREFERENCES

Utility function:  $U [c_t^i(s^t)] = \frac{c_t^i(s^t)^{1-\gamma}}{1-\gamma}$ ,  $\gamma > 0$ .

Maximizing the expected sum of discounted utilities over all  $t, s^t$  s.t. the individual budget constraint, gives the following equilibrium condition for agents  $i$  and 1 (an application of (8.5.5)):

$$c_t^i(s^t) = c_t^1(s^t) \left( \frac{\mu_i}{\mu_1} \right)^{-\frac{1}{\gamma}}.$$

Summing over  $i$ ,

$$\sum_i c_t^i(s^t) = (\mu_1)^{\frac{1}{\gamma}} c_t^1(s^t) \sum_i (\mu_i)^{-\frac{1}{\gamma}} = \sum_i y_t^i(s^t),$$

and so

$$c_t^1(s^t) = (\mu_1)^{-\frac{1}{\gamma}} \left[ \sum_i (\mu_i)^{-\frac{1}{\gamma}} \right]^{-1} \sum_i y_t^i(s^t) = \text{const} \sum_i y_t^i(s^t).$$

Thus, individual consumption at  $t, s^t$  is *perfectly* correlated with the *aggregate endowment* and *aggregate consumption*.

## RISK SHARING. SUMMARY

- Consumption of agents  $i$  and  $j$  are constant fractions of one another, independent of  $t, s^t$ .
- Conditional on time  $t$ , history  $s^t$ ,  $c_t^i(s^t)$  is independent of the individual endowment at  $t, s^t$ , for all  $i$ .
- There is an extensive cross-history, cross-time consumption smoothing. Constant-fraction consumption allocation comes from: 1) the complete markets assumption; 2) a homothetic utility function.

## EXAMPLE (8.6.2): NO AGGREGATE UNCERTAINTY

Assume  $s_t \in [0, 1]$ ;  $s^t$  is a stochastic event;  $i = 1, 2$ ;  $y_t^1(s^t) = s_t$ ;  
 $y_t^2(s^t) = 1 - s_t$ .

Aggregate endowment at  $t, s^t$ :

$$\sum_i y_t^i(s^t) = s_t + (1 - s_t) = 1.$$

Thus, it immediately follows from (8.5.7) that  $c_t^1(s^t)$  is constant for each  $t, s^t$ .  $c_t^2(s^t) = 1 - c_t^1(s^t)$  is also constant  $\forall t, s^t$ .

Equilibrium allocation:

$$c_t^i(s^t) = \bar{c}^i.$$

From (8.5.4) it follows

$$\beta^t \pi_t(s^t) \frac{u'(\bar{c}^i)}{\mu_i} = q_t^0(s^t), \quad \forall i, t, s^t \quad (8.6.2)$$

Agent  $i$ 's b.c.:

$$\sum_t \sum_{s^t} q_t^0(s^t) [\bar{c}^i - y_t^i(s^t)] = 0,$$

or

$$\bar{c}^i \sum_t \sum_{s^t} q_t^0(s^t) - \sum_t \sum_{s^t} q_t^0(s^t) y_t^i(s^t) = 0.$$

Plug the definition for  $q_t^0(s^t)$  from (8.5.4), to obtain:

$$\begin{aligned}
\bar{c}^i \frac{u'(\bar{c}^i)}{\mu_i} \sum_t \sum_{s^t} \pi_t(s^t) \beta^t - \frac{u'(\bar{c}^i)}{\mu_i} \sum_t \sum_{s^t} \pi_t(s^t) \beta^t y_t^i(s^t) &= 0 \\
\bar{c}^i \sum_t \beta^t \sum_{s^t} \pi_t(s^t) - \sum_t \sum_{s^t} \pi_t(s^t) \beta^t y_t^i(s^t) &= 0 \\
c^i = (1 - \beta) \sum_t \sum_{s^t} \pi_t(s^t) \beta^t y_t^i(s^t) & \quad (8.6.3)
\end{aligned}$$

Is this allocation feasible? For each  $t, s^t$ :

$$\begin{aligned}
\sum_i \bar{c}^i &= (1 - \beta) \sum_i \sum_t \sum_{s^t} \pi_t(s^t) \beta^t y_t^i(s^t) \\
&= (1 - \beta) \sum_t \sum_{s^t} \sum_i \pi_t(s^t) \beta^t y_t^i(s^t) \\
&= (1 - \beta) \sum_t \sum_{s^t} \pi_t(s^t) \beta^t * 1 \\
&= (1 - \beta) \sum_t \beta^t \sum_{s^t} \pi_t(s^t) \\
&= (1 - \beta) / (1 - \beta) = 1.
\end{aligned}$$

In this example, consumer perfectly smoothes consumption over time and across histories since aggregate endowment does not fluctuate across time.

## EXAMPLE: PERIODIC ENDOWMENT PROCESSES

$s_t \in [0, 1]$ ,  $s_t$  deterministic. Economy starts at  $t = 0$ .  
Agent 1's endowment sequence: 1, 0, 1, 0, 1, ... Agent 2's  
endowment sequence: 0, 1, 0, 1, 0, ...  $\pi_t(s^t) = 1$  since  $s^t$  is  
deterministic.

(8.5.4) implies:

$$q_t^0(s^t) = \beta^t \pi_t(s^t) \frac{u'(\bar{c}^i)}{\mu_i} = \beta^t \frac{u'(\bar{c}^i)}{\mu_i}.$$

Normalize  $q_0^0(s_0) = 1$ , i.e., set the time-0 good as a numeraire.  
Thus,  $1 = \beta^0 \frac{u'(\bar{c}^i)}{\mu_i} = \frac{u'(\bar{c}^i)}{\mu_i}$ . It follows that

$$q_t(s^t) = \beta^t.$$

(8.6.3) implies:

$$\begin{aligned}\bar{c}^i &= (1 - \beta) \sum_t \sum_{s^t} \pi_t(s^t) \beta^t y_t^i(s^t) \\ &= (1 - \beta) \sum_t \sum_{s^t} \beta^t y_t^i(s^t).\end{aligned}$$

$$\begin{aligned}\bar{c}^1 &= (1 - \beta)(1 + \beta * 0 + 1 * \beta^2 + 0 * \beta^3 + 1 * \beta^4 + \dots) = \\ &= (1 - \beta) \frac{1}{1 - \beta^2} = \frac{1}{1 + \beta}.\end{aligned}$$

$$\bar{c}^2 = 1 - \bar{c}^1 = \frac{\beta}{1 + \beta}.$$

Note that a lifetime richer agent 1 consumes more.

# Asset pricing in complete markets

In complete markets, an asset will offer a bundle of history-contingent claims, each component of which has been already priced in the market. In this case, the asset is viewed as *redundant*.

## Pricing redundant assets

An asset brings  $\{d_t(s^t)\}_{t=0}^{\infty}$ —a stream of claims on  $t, s^t$  consumption (in physical units). (Think of dividends here.)

The time 0 price of this asset in terms of time 0, history  $s_0$  consumption must be:

$$p_0^0(s_0) = \sum_{t=0}^{\infty} \sum_{s^t} d_t(s^t) q_t^0(s^t).$$

Why? We can replicate this asset by purchasing consumption bundles  $\{d_t(s^t)\}_{t=0}^{\infty}$  at the total cost  $\sum_{t=0}^{\infty} \sum_{s^t} d_t(s^t) q_t^0(s^t)$ . Hence, the asset must cost  $p_0^0(s_0)$ .

A *riskless consol* brings 1 unit of consumption in *each*  $t$ , regardless of  $s^t$ .

To replicate its payoffs, we need to purchase  $q_t^0(s^t)$  units of claims for each  $t, s^t$ . The price of this asset, therefore, is:

$$\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t).$$

A *riskless strip* pays  $d_\tau = 1$  if  $\tau = t \geq 0$ , 0—otherwise. Thus, the strips pay off only at time  $\tau$ , regardless of the history realization  $s^\tau$ .

To replicate the asset's payoffs, we need to purchase, at  $t = 0$ ,  $q_\tau^0(s^\tau)$  units of consumption for time  $\tau$ , all histories  $s^\tau$ . Thus, the price of the strip must be

$$\sum_{s^\tau} q_\tau^0(s^\tau).$$

## Pricing tail assets

*Tail assets* pay  $\{d_t(s^t)\}_{t \geq \tau, \tau \geq 1}^\infty$ . At time 0, this asset will cost

$$p_\tau^0(s^\tau) = \sum_{s^t | s^\tau} \sum_{t \geq \tau} d_t(s^t) q_t^0(s^t).$$

Units of this price are in terms of time 0, state  $s^0 = s_0$  goods.

Quote the price in *units of time  $\tau$ , state  $s^\tau$  goods*. Divide the price by  $q_\tau^0(s^\tau)$ , to obtain:

$$p_\tau^\tau(s^\tau) \equiv \frac{p_\tau^0(s^\tau)}{q_\tau^0(s^\tau)} = \sum_{t \geq \tau} \sum_{s^t | s^\tau} d_t(s^t) \frac{q_t^0(s^t)}{q_0^\tau(s^\tau)}. \quad (8.7.4)$$

- $p_\tau^0(s^\tau)$  measures the amount of goods at time 0 state  $s^0$  it costs to purchase the asset (# of time 0 goods per 1 asset).
- $q_\tau^0(s^\tau)$  measures the amount of goods at time 0 state  $s^0$  it costs to purchase the claim to one good at time  $\tau$  state  $s^\tau$  (# of time 0 goods per 1 unit of goods of time  $\tau$  state  $s^\tau$ ).
- It follows that  $\frac{p_\tau^0(s^\tau)}{q_\tau^0(s^\tau)}$  measures the amount of goods of time  $\tau$  state  $s^\tau$  it costs to purchase the asset.

Define  $p_\tau^\tau(s^\tau) = \frac{p_\tau^0(s^\tau)}{q_\tau^0(s^\tau)}$ —the price of the tail asset in terms of the time  $\tau$ , history  $s^\tau$  consumption goods.

Note that in (8.7.4)

$$\begin{aligned} q_t^\tau(s^t) &\equiv \frac{q_t^0(s^t)}{q_\tau^0(s^\tau)} = \frac{\mu_i^{-1} \beta^t u' [c_t^i(s^t)] \pi_t(s^t)}{\mu_i^{-1} \beta^\tau u' [c_\tau^i(s^\tau)] \pi_\tau(s^\tau)} \\ &= \beta^{t-\tau} \frac{u' [c_t^i(s^t)] \pi_t(s^t)}{u' [c_\tau^i(s^\tau)] \pi_\tau(s^\tau)}. \end{aligned} \quad (8.7.5)$$

Since the LHS of the equation is *not* individual-dependent, the RHS should not be too. Note that  $\frac{\pi_t(s^t)}{\pi_\tau(s^\tau)} = \frac{\pi(s^t|s^\tau)\pi(s^\tau)}{\pi(s^\tau)}$ .

Thus,

$$p_\tau^\tau(s^\tau) = \sum_{t \geq \tau} \sum_{s^t | s^\tau} \left[ \beta^{t-\tau} \frac{u' [c_t^i(s^t)]}{u' [c_\tau^i(s^\tau)]} \pi_t(s^t | s^\tau) \right] d_t(s^t). \quad (8.7.6)$$

## Pricing one-period returns

Replace  $t$  with  $\tau + 1$  in (8.7.5), to obtain:

$$q_{\tau+1}^{\tau}(s^{\tau}) = \beta \frac{u'[c_{\tau+1}^i(s^{\tau+1})]}{u'[c_{\tau}^i(s^{\tau})]} \pi_{\tau+1}(s^{\tau+1} | s^{\tau}).$$

If we want to find the price at time  $\tau$ , history  $s^{\tau}$  of a claim to a random payoff  $\omega(s_{\tau+1})$ , we'd use the following formula:

$$p_{\tau}^{\tau}(s^{\tau}) = \sum_{s_{\tau+1}} q_{\tau+1}^{\tau}(s^{\tau+1}) \omega(s_{\tau+1}) = E_{\tau} \left[ \beta \frac{u'[c_{\tau+1}]}{u'[c_{\tau}]} \omega(s_{\tau+1}) \right]. \quad (8.7.7)$$

Let the one-period gross return on the asset with payoff  $\omega(s_{\tau+1})$  be  $R_{\tau+1} \equiv \frac{\omega(s_{\tau+1})}{p_{\tau}^T(s^{\tau})}$ .

Then (8.7.7) becomes:

$$E_{\tau} \left[ \beta \frac{u' [c_{\tau+1}]}{u' [c_{\tau}]} R_{\tau+1} \right] = E_{\tau} [m_{\tau+1} R_{\tau+1}] = 1,$$

where  $m_{\tau+1} = \beta \frac{u' [c_{\tau+1}]}{u' [c_{\tau}]}$  is called a **stochastic discount factor**.

## Sequential trading: Arrow securities

- At each date  $t \geq 0$ , trades occur in a set of claims to one-period-ahead state-contingent consumption.
- With a full array of one-period-ahead claims, the sequential trading arrangement attains the same competitive allocation as we described earlier.

- If trading takes place in period 0, what is the implied wealth of household  $i$  at time  $t$  after history  $s^t$ ? It is the value of household's purchased claims net of its outstanding liabilities.
- When history  $s^t$  is realized discard all claims and liabilities conditional on another initial history.
- Household's net claim of goods at  $\tau$  is  $c_\tau(s^\tau) - y_\tau(s^\tau)$ .
- Household's wealth in terms of the date  $t$ , history  $s^t$  consumption good is

$$\Upsilon_t^i(s^t) = \sum_{\tau=t}^{\infty} \sum_{s^\tau | s^t} q_\tau^t(s^\tau) [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)]. \quad (8.8.1)$$

Notice that

$$\sum_{i=1}^I \Upsilon_t^i(s^t) = \sum_{\tau=t}^{\infty} \sum_{s^\tau | s^t} q_\tau^t(s^\tau) \underbrace{\sum_{i=1}^I [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)]}_{=0 \text{ by feasibility}}, \text{ for each } t, s^t.$$

We can match up the time  $t$ , history  $s^t$  wealth of the household in the sequential economy with the tail wealth  $\Upsilon_t^i(s^t)$  from the Arrow-Debreu economy as in equation (8.8.1).

## Debt limits

“Natural debt limits:” it should be feasible for the consumer to repay his state-contingent debt in every possible state.

Let  $q_\tau^t(s^\tau)$  be the price of the good at time  $\tau$ , history  $s^\tau$  in units of the good at time  $t$ , history  $s^t$ . The value of the tail of the endowment sequence at time  $t$  is

$$A_t^i(s^t) = \sum_{\tau=t}^{\infty} \sum_{s^\tau | s^t} q_\tau^t(s^\tau) y_\tau^i(s^\tau). \quad (8.8.2)$$

$A_t^i(s^t)$  is the natural debt limit at time  $t$  and history  $s^t$ .

At each date  $t \geq 0$ , households trade claims to date  $t$  consumption, whose payment is contingent on the realization of  $s_{t+1}$ .

- $\tilde{a}_t^i(s^t)$ —the claims to time  $t$  consumption that household brings to time  $t$ , history  $s^t$ .
- $\tilde{Q}_t(s_{t+1}|s^t)$ —the price of one unit of time  $t + 1$  consumption contingent on realization of  $s_{t+1}$  at  $t + 1$ , given the history at  $t$  is  $s^t$ .
- The time  $t$  history  $s^t$  budget constraint is

$$\underbrace{\tilde{c}_t^i(s^t) + \sum_{s_{t+1}} \tilde{a}_{t+1}^i(s_{t+1}, s^t) \tilde{Q}_t(s_{t+1}, s^t)}_{\text{time } t \text{ history } s^t \text{ expenditures}} \leq \underbrace{y_t^i(s^t) + \tilde{a}_t^i(s^t)}_{\text{time } t \text{ history } s^t \text{ resources}}$$

(8.8.3)

## State-by-state borrowing constraints

$$\tilde{a}_{t+1}^i(s^{t+1}) \geq -A_{t+1}^i(s^{t+1}). \quad (8.8.4)$$

The Lagrangian is

$$\begin{aligned} L^i = & \sum_{t=0}^{\infty} \sum_{s^t} \{ \beta^t u(\tilde{c}_t^i(s^t)) \pi_t(s^t) \\ & + \eta_t^i(s^t) \left[ y_t^i(s^t) + \tilde{a}_t^i(s^t) - \tilde{c}_t^i(s^t) - \sum_{s_{t+1}} \tilde{a}_{t+1}^i(s_{t+1}, s^t) \tilde{Q}_t(s_{t+1}|s^t) \right] \\ & + \sum_{s_{t+1}} \nu_t^i(s^t; s_{t+1}) [A_{t+1}^i(s_{t+1}, s^t) + \tilde{a}_{t+1}^i(s_{t+1}, s^t)] \}, \end{aligned}$$

for a given  $\tilde{a}_0^i(s_0)$ .

FOCs w.r.t.  $\tilde{c}_t^i(s^t)$  and  $\{\tilde{a}_{t+1}^i(s_{t+1}, s^t)\}_{s_{t+1}}$  are

$$\beta^t u' [\tilde{c}_t^i(s^t)] \pi_t(s^t) - \eta_t^i(s^t) = 0 \quad (8.8.5a)$$

$$-\eta_t^i(s^t) \tilde{Q}_t(s_{t+1}|s^t) + \nu_t^i(s^t; s_{t+1}) + \eta_{t+1}^i(s_{t+1}, s^t) = 0, \quad (8.8.5b)$$

for all  $t$ ,  $s^t$ , and  $s_{t+1}$ . At the optimum, natural borrowing constraints are not binding, and so  $\nu_t^i(s^t; s_{t+1}) = 0$ . Thus,

$$\tilde{Q}_t(s_{t+1}|s^t) = \beta \frac{u' [\tilde{c}_{t+1}^i(s^{t+1})]}{u' [\tilde{c}_t^i(s^t)]} \pi_t(s^{t+1}|s^t), \quad (8.8.6)$$

for all  $t$ ,  $s^t$ , and  $s_{t+1}$ .

## Equivalence of allocations

AD equilibrium is also a sequential-trading equilibrium with a particular initial distribution of wealth. Make a guess that

$$\begin{aligned} q_{t+1}^0(s^{t+1}) &= \tilde{Q}_t(s_{t+1}|s^t)q_t^0(s^t) \\ \tilde{Q}_t(s_{t+1}|s^t) &= \frac{q_{t+1}^0(s^{t+1})}{q_t^0(s^t)} = q_{t+1}^t(s^{t+1}). \end{aligned} \quad (8.8.7)$$

For the Arrow-Debreu economy, the following Euler equation should hold

$$\frac{\beta u' [c_{t+1}^i(s^{t+1})] \pi (s^{t+1}|s^t)}{u' [c_t^i(s^t)]} = \frac{q_{t+1}^0(s^{t+1})}{q_t^0(s^t)} = \tilde{Q}_t(s_{t+1}|s^t) \quad (8.8.8)$$

The equation is equivalent with the sequential-trading FOC (8.8.6). We need to establish that  $\tilde{c}_t^i(s^t) = c_t^i(s^t)$ .

Choose the initial wealth in the sequential-trading equilibrium so that this equilibrium duplicates the Arrow-Debreu equilibrium. Let  $\tilde{a}_0^i(s_0)$  for all  $i$ —households rely on their endowment stream to finance its consumption.

Need to show that household  $i$  will be able to finance  $\{c_t^i(s^t)\}$  and there will be no room to improve on that allocation.

Guess that

$$\tilde{a}_{t+1}^i(s_{t+1}, s^t) = \Upsilon_{t+1}^i(s^{t+1}), \text{ for all } s_{t+1}.$$

The value of this portfolio is

$$\begin{aligned} \sum_{s_{t+1}} \tilde{a}_{t+1}^i(s_{t+1}, s^t) \tilde{Q}_t(s_{t+1}|s^t) &= \sum_{s^{t+1}|s^t} \Upsilon_{t+1}^i(s^{t+1}) q_{t+1}^t(s^{t+1}) \\ &= \sum_{\tau=t+1}^{\infty} \sum_{s^\tau|s^{t+1}} \sum_{s^{t+1}|s^t} q_{t+1}^t(s^{t+1}) q_\tau^{t+1}(s^\tau) [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)] \\ &= \sum_{\tau=t+1}^{\infty} \sum_{s^\tau|s^t} q_\tau^t(s^\tau) [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)]. \end{aligned} \tag{8.8.9}$$

At time 0, the budget constraint (8.8.3) implies

$$\tilde{c}_0(s_0) + \sum_{t=1}^{\infty} \sum_{s^t} q_t^0(s^t) [c_t^i(s^t) - y_t^i(s^t)] = y_0^i(s_0) + \underbrace{\tilde{a}_0^i(s_0)}_{=0}$$

$$\tilde{c}_0(s_0) + \sum_{t=1}^{\infty} \sum_{s^t} q_t^0(s^t) [c_t^i(s^t) - y_t^i(s^t)] = y_0^i(s_0) - c_0^i(s_0) + c_0^i(s_0)$$

$$\tilde{c}_0(s_0) + \underbrace{\sum_{t=0}^{\infty} \sum_{s^t} q_t^0(s^t) [c_t^i(s^t) - y_t^i(s^t)]}_{=0 \text{ since } u' > 0} = c_0^i(s_0).$$

Thus, the proposed portfolio is affordable and in period 0 and the associated consumption level is the same as in the Arrow-Debreu competitive equilibrium.

From (8.8.9),

$$\begin{aligned}\sum_{s_{t+1}} \tilde{a}_{t+1}^i(s_{t+1}, s^t) \tilde{Q}_t(s_{t+1} | s^t) &= \sum_{\tau=t}^{\infty} \sum_{s^\tau} q_\tau^t(s^\tau) [c_\tau^i(s^\tau) - y_\tau^i(s^\tau)] \\ &\quad - \underbrace{q_t^t(s^t)}_{=1} [c_t^i(s^t) - y_t^i(s^t)] \\ &= \Upsilon_t^i(s^t) - [c_t^i(s^t) - y_t^i(s^t)].\end{aligned}$$

From (8.8.3) it follows,

$$\begin{aligned}\tilde{c}_t^i(s^t) + \Upsilon_t^i(s^t) - [c_t^i(s^t) - y_t^i(s^t)] &= y_t^i(s^t) + \tilde{a}_t^i(s^t) \\ \tilde{c}_t^i(s^t) + \Upsilon_t^i(s^t) &= c_t^i(s^t) + \tilde{a}_t^i(s^t).\end{aligned}$$

Thus, when  $\tilde{a}_t^i(s^t) = \Upsilon_t^i(s^t)$ ,  $\tilde{c}_t^i(s^t) = c_t^i(s^t)$ , for all  $t, s^t$ .